

A Realistic Model for Rescue Operations after an Earthquake^{*} (Extended Abstract)

Tiziana Calamoneri and Federico Corò

Department of Computer Science
Sapienza University of Rome, Italy

When natural disasters, such as earthquakes, occur, rapid deployment of rescue operations is extremely important, in order to avoid an increase in loss of lives. Moreover, we cannot even count on the presence of public communications networks and infrastructures such as internet or mobile communication, that could be disrupted.

In this context, *Unmanned Aerial Vehicles* (UAVs) can make the difference, contributing to reduce the costs, resources, and human risks in support of more efficient Search & Rescue operations.

Recently, several papers have been published proposing many theoretical models for handling algorithmic optimization problems related to the effective rescue operations of UAVs in case of natural disasters.

Nevertheless, since the problems related to these issues are very complex, most of such models appear too distant from real life and assume some hypotheses that are not reasonable in practice. For example, in the majority of papers, UAVs are supposed not to have battery constraints (e.g., [2, 8]) that are instead very pressing; some authors assume that there exists a functioning cellular network to exploit the use of mobile phones or to send stream videos from the UAVs to the base [1–3, 12]; finally, the costs that a UAV will face during its trip are usually considered as fully known [5, 11], while, in real-life scenarios, a UAV could spend a different amount of time from the expected one to explore or traverse an area.

In this work, we introduce some realistic constraints to the theoretical model. Namely, first of all, we do not require UAVs to use internet to communicate with the base.

Secondly, the battery constraint is considered as very pressing (so we do not allow any elasticity).

Then, we assign a priority to every site, to prioritize the most critical buildings, e.g., hospitals, and schools.

Finally, we introduce uncertainty in the overflight times, due to the autonomous decision of a UAV to deepen the overflight of a site when a possible survivor is detected.

^{*} Copyright © 2020 for this paper by its authors. Use permitted under Creative Commons License Attribution 4.0 International (CC BY 4.0).

It is reasonable to assume that civil defense constantly possesses an updated map of every town and village considered at risk and, for every relevant building, it has the information about the necessary time to overfly it entirely and the level of importance. Civil defense can have in very few minutes a clear idea of the boundary of the affected area and set the level of importance of the buildings and sites that have been subjected to some damages based on emergency calls and on the moment when the earthquake happened: just to make a single example, if it was during the night or in summer, schools are assigned a very low priority; vice-versa, if it happened in the morning of the school year, schools will receive the highest priority. Civil Defense usually has a “National Plan for seismic risk prevention” in which each region of the country is analyzed, to assess the risk, that is, to identify the areas that, in the event of a seismic wave, may be subject to amplification phenomena and provide useful indications for urban planning and rescue operation.

Then, we consider that on hand of civil defense is a fleet of q homogeneous asynchronous UAVs u_1, \dots, u_q taking off from the operations center v_0 . Each UAV has a battery that corresponds to $B \in \mathbb{R}^+$ units of flying time; when it is run-down, it takes $R \in \mathbb{R}^+$ (with $R \geq B$: current estimate is, for example, $1.5B \leq R \leq 2B$ [6]) time units to be fully recharged and this can only be done at v_0 .

We consider the emergency situation after a natural disaster (e.g., an earthquake) and the main issue of overflying the whole zone employing the fleet of UAVs, acquiring information as soon as possible by sensors and/or cameras mounted on UAVs, to have a clear idea about the places where rescue teams are more necessary. So, the aim of our problem consists of assigning to every UAV an ordered set of sites to overfly so that the whole fleet can inspect all the areas of interest as soon as possible giving precedence to the sites with higher priority. Since our situation is time-critical, we allow having a certain number $b \geq 0$ of batteries that can be separately recharged while UAVs are in flight. These batteries can be quickly swapped out in seconds with those exhausted without causing any delay using methods such as those proposed in [10, 9]. We theoretically analyze each possible value of b and associate to it the corresponding value of the completion time of rescue operations, so that civil defense organizers can have an immediate idea about the necessary time to conclude operations, knowing all the parameters. In particular, we derive that $\lceil \frac{R}{B} \rceil \cdot q$ is the minimum number of batteries b necessary to guarantee no additional time in which UAVs remain inactive waiting for battery recharging.

We consider two different scenarios, according to the devices civil defense can count on and the assumptions we are able to do on the environment.

In the *first scenario*, each UAV is equipped with an RGB commercial camera, it has neither computational power nor communication devices. In this case, the only thing UAVs can do is to follow an assigned route through some sensible targets, record a video and bring it back to the operations center; there, an image detection/recognition tool will be able to detect in parallel on all videos possible survivors needing help in real-time. The fact that the videos must be brought

back to the base before being processed can cause a waste of time. However, the positive side of this scenario is that the fleet can be easily constituted by a large number of UAVs because they are very cheap.

In the *second scenario*, we considered high-cost UAVs, so they have a processing unit with good computational power, and can be equipped with a device allowing them to communicate with the base at every time (e.g., a radio); so we can assume they are able to scan in real-time the videos while recording them and immediately detect and recognize whether there is an emergency; in case some people requesting help are detected, they promptly decrease their altitude, closer to the possible survivors, acquire additional information concerning detected people possibly in difficulty and send to the base a radio message, so that a rescue team is sent. Note that the radio onboard is also exploited to inform the base in real-time about the battery status and the capability to overfly a further site.

In order to model the described context, we consider a complete node- and edge-weighted graph $G = (V, E, dist, \sigma, p)$, where:

- $V = \{v_0, v_1, \dots, v_n\}$: v_1, \dots, v_n are the sites representing the points of interest inside the area, i.e., buildings (or areas) that need to be explored, while v_0 represents the operations centre.
- Each node is associated with two different values, a function p that represents the priority and a function σ , that represents the time needed to explore a site. In the first scenario, $\sigma(v)$ represents the only (fixed) time to fly over each site. In the second scenario, it may be necessary to add to $\sigma(v)$ a second contribution due to the time that the UAV needs to take more precise information concerning detected people and send a message to the base.
- Every (undirected) edge $e = (v_i, v_j) \in E$, with $i \neq j$, is associated with a *distance* function $dist$ that represents the symmetric traveling distance between two nodes.

The aim of our problem consists of assigning to every UAV an ordered set of sites to overfly so that the whole fleet can inspect all the areas of interest as soon as possible giving precedence to the sites with higher priority. Since UAVs are prone to their battery constrain, they could be not able to visit all sites assigned to them in a unique flight, but need to perform more than one flight, after each of them returning to the base for either recharging or substituting its battery. Thus, we have to assign to each UAV an ordered set of *cycles* all passing through v_0 , whose weights are bounded by B , such that all nodes of the graph are overflown, giving precedence to those with higher priority, and in such a way an opportune optimization function is minimized. Of course, in the two scenarios, the situation is different and hence also the solution is. In particular, it is worth to be noticed that, while in the first scenario the times needed to explore each site are exact, in the second one they are not.

Let us define a *sequence* \mathcal{S}_{u_i} as a set of cycles that are assigned to UAV u_i . Then, a solution set *SOL* is set of sequences for each UAV that covers V , i.e., if for every $v \in V$ there exists a cycle in *SOL* containing v .

We informally define the completion time of a site as the time necessary to know whether there are people needing help in it. Given a node, in the first scenario, this occurs only after that the video of the cycle including it has been delivered to the base and the portion corresponding to it has been analyzed. In the second scenario, instead, this occurs as soon as the node has been completely overflown. In this extended abstract, we omit the formalization of the definition of *completion time of site v* , $cost^{(i)}(v, SOL)$, where $i = I, II$ and refer to, respectively, to the first and second scenario described above.

We now define two functions that exploit the concept of completion time of a site and are useful to evaluate the goodness of a solution: The *weighted latency*, wL and the *completion time ct* . Formally,

$$wL^{(i)}(SOL) = \frac{1}{n} \sum_{v \in V} p(v) cost^{(i)}(v, SOL) \text{ where } i = I, II. \quad (1)$$

and

$$ct^{(i)}(SOL) = \max_{v \in V} cost^{(i)}(v, SOL) \text{ where } i = I, II.$$

Observe that, while $ct^{(i)}$ does not take into account the priority of nodes, not only in $wL^{(i)}$ priorities are taken into account, but sites with the highest priority should be served first to have a smaller value of $wL^{(i)}$, in order to associate a larger multiplicative factor (corresponding to priority) to a smaller sum of costs. We call *HPSF (Higher Priority Sites First)* this implicit problem requirement.

Clearly better solutions for our scenario will have smaller values of ct and wL , from which it follows the formal definition of our problem:

Given a complete node- and edge-weighted graph $G = (V, E, dist, \sigma, p)$ and a set of q homogeneous UAVs u_1, \dots, u_q , the problem Cover by Multitrips with Priorities (CMP) consists in finding a set SOL of q sequences such that $ct^{(i)}(SOL)$, where $i = I, II$, is minimized and Requirement HPSF is addressed.

It is easy to see that CMP is a generalization of the well known m -TSP [4] (achievable with $B = \infty$, and, for each node $v \in V$, $p(v) = 1$ and full-knowledge of function $\sigma(v)$, i.e., either the second contribution of function $\sigma(v)$ is zero or known in advance). Thus, in turn, it is strongly *NP*-hard. It seems instead not possible to inherit any approximability result due to the presence of B , as pointed out in [7] for a different problem, without relaxing the battery constraint (and accepting, for instances, cycles with completion time $B + \epsilon$ for some $\epsilon > 0$), that in our context is inadmissible.

We propose four different heuristics, each one based on a different consideration but all based on a meta-algorithm that is run at the operations center; it consists of a number of iterations going on until all sites have been overflown; at each iteration, a cycle for each UAV is produced as output and its length is guaranteed to be bounded by B . The meta-algorithm (and hence all the heuristics) provably terminate in a finite time whose an upper bound can be analytically determined in both scenarios.

This meta-algorithm can be run also as a preprocessing algorithm in order to do some preliminary estimates on the necessary completion time and the best number of necessary additional batteries and of UAVs.

We evaluate and compare the performance of the proposed heuristics running extended experimentation on two types of randomly generated graphs: we position n points, varying n between 50 and 200, placed either uniformly at random or with a Poisson point process in a unit square. This latter distribution is intended to simulate building collapses in some specific zones.

We remand to the full version of the paper the interested reader for a detailed description of the graphics out-coming from the experimental results.

References

1. Albanese, A., Sciancalepore, V., Costa-Pérez, X.: SARDO: an automated search-and-rescue drone-based solution for victims localization. *CoRR* **abs/2003.05819** (2020)
2. Baker, C.A.B., Ramchurn, S.D., Teacy, W.T.L., Jennings, N.R.: Planning search and rescue missions for UAV teams. In: In 22nd European Conference on Artificial Intelligence (ECAI). pp. 1777–1782 (2016)
3. Bejiga, M.B., Zegada, A., Nouffidj, A., Melgani, F.: A convolutional neural network approach for assisting avalanche search and rescue operations with UAV imagery. *Remote Sensing* **9**(2), 100 (2017)
4. Bektas, T.: The multiple traveling salesman problem: an overview of formulations and solution procedures. *omega* **34**(3), 209–219 (2006)
5. Bevacqua, G., Cacace, J., Finzi, A., Lippiello, V.: Mixed-initiative planning and execution for multiple drones in search and rescue missions. In: Proceedings of the Twenty-Fifth International Conference on Automated Planning and Scheduling, ICAPS. pp. 315–323 (2015)
6. Drones, A.: Airborne drones (2020), <https://www.airboredrones.co/natural-disasters/>
7. Jothi, R., Raghavachari, B.: Approximating the k-traveling repairman problem with repair times. *J. Discrete Algorithms* **5**(2), 293–303 (2007)
8. Kim, D., Xue, L., Li, D., Zhu, Y., Wang, W., Tokuta, A.O.: On theoretical trajectory planning of multiple drones to minimize latency in search-and-reconnaissance operations. *IEEE Trans. Mob. Comput.* **16**(11), 3156–3166 (2017)
9. Liu, Z.N., Liu, X.Q., Yang, L.J., Leo, D., Zhao, H.: An autonomous dock and battery swapping system for multirotor uav. Unpublished, <https://www.researchgate.net/publication/325077351> **10** (2018)
10. Liu, Z., hao Wang, Z., Leo, D., Liu, X., Zhao, H.: QUADO: an autonomous recharge system for quadcopter. In: 2017 IEEE International Conference on Cybernetics and Intelligent Systems (CIS) and IEEE Conference on Robotics, Automation and Mechatronics (RAM). pp. 7–12 (2017)
11. Modares, J., Ghanei, F., Mastronarde, N., Dantu, K.: UB-ANC planner: Energy efficient coverage path planning with multiple drones. In: International Conference on Robotics and Automation, ICRA. pp. 6182–6189 (2017)
12. Sun, J., Li, B., Jiang, Y., Wen, C.: A camera-based target detection and positioning UAV system for search and rescue (SAR) purposes. *Sensors* **16**(11), 1778 (2016)